

Better Orders for Saturated Cost Partitioning in Optimal Classical Planning

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University of Basel

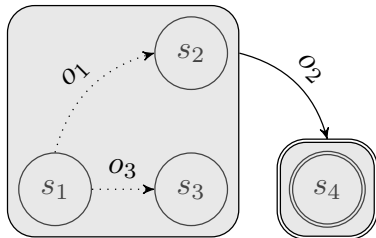
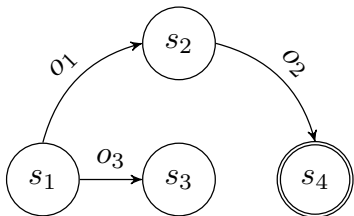
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Setting

- optimal classical planning
- A* search + admissible heuristic
- abstraction heuristics

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Problem

- single heuristic unable to capture enough information

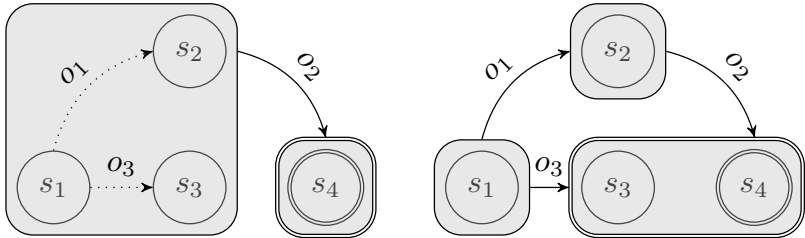
Problem

- single heuristic unable to capture enough information
→ use **multiple heuristics**

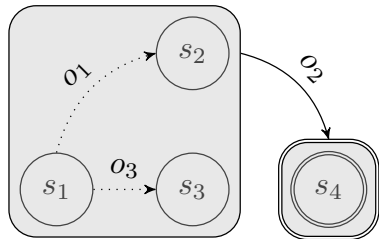
Problem

- single heuristic unable to capture enough information
→ use **multiple heuristics**
- how to **combine** multiple heuristics admissibly?

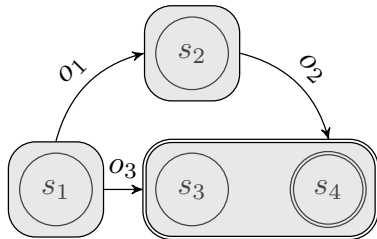
Multiple Heuristics



Multiple Heuristics

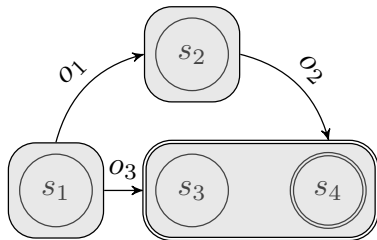
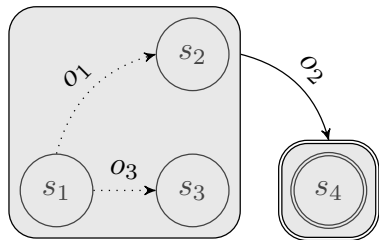


$$h_1(s_1) = 1$$



$$h_2(s_1) = 1$$

Multiple Heuristics



- maximizing only **selects** best heuristic $\rightarrow h(s_1) = 1$

Multiple Heuristics: Cost Partitioning

Cost Partitioning

- split operator costs among heuristics
- total costs must not exceed original costs

→ combines heuristics

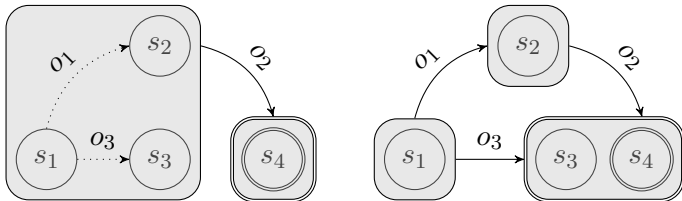
→ allows summing heuristic values admissibly

Saturated Cost Partitioning

Seipp & Helmert (ICAPS 2014)

Saturated Cost Partitioning Algorithm

- order heuristics
- for each heuristic h :
 - use minimum costs preserving all estimates of h
 - use remaining costs for subsequent heuristics

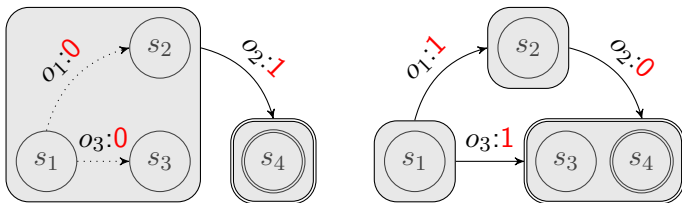


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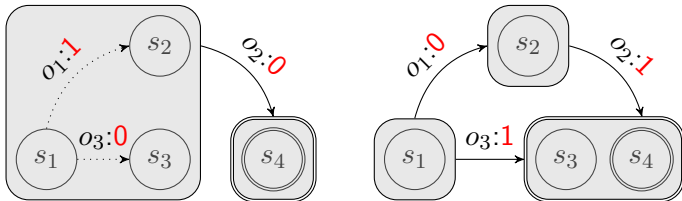
$$h_{\langle h_1, h_2 \rangle}^{\text{SCP}}(s_1) = 1 + 1 = 2$$

Saturated Cost Partitioning

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Saturated Cost Partitioning Algorithm

- **order** heuristics
- for each heuristic h :
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$$h_{\langle h_1, h_2 \rangle}^{\text{SCP}}(s_1) = 1 + 1 = 2$$

$$h_{\langle h_2, h_1 \rangle}^{\text{SCP}}(s_1) = 1 + 0 = 1$$

Orders for Saturated Cost Partitioning

Seipp, Keller & Helmert (AAAI 2017)

- orders can be arbitrarily bad
→ hill climbing in space of orders
- a single order might only be good for a single state
→ use multiple orders
- using too many orders slows down evaluation
→ use diverse orders

Combining Heterogeneous Abstraction Heuristics

ICAPS 2017

$h_{\text{HC}}^{\text{SCP}}$ $h_{\text{Sys}}^{\text{SCP}}$ $h_{\text{Cart}}^{\text{SCP}}$

Tasks (1667) 805 852 965

Combining Heterogeneous Abstraction Heuristics

ICAPS 2017

$h_{\text{HC}}^{\text{SCP}}$ $h_{\text{Sys}}^{\text{SCP}}$ $h_{\text{Cart}}^{\text{SCP}}$ $h_{\text{HC, Sys, Cart}}^{\text{SCP}}$

Tasks (1667) 805 852 965 1006

Combining Heterogeneous Abstraction Heuristics

ICAPS 2017

$$\begin{array}{c} \hline h_{\text{HC}}^{\text{SCP}} \quad h_{\text{Sys}}^{\text{SCP}} \quad h_{\text{Cart}}^{\text{SCP}} \\ \hline h_{\text{HC, Sys, Cart}}^{\text{SCP}} \quad \max(h_{\text{HC}}^{\text{SCP}}, h_{\text{Sys}}^{\text{SCP}}, h_{\text{Cart}}^{\text{SCP}}) \end{array}$$

Tasks (1667)	805	852	965	1006	1032
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Combining Heterogeneous Abstraction Heuristics

ICAPS 2017

$$\frac{h_{\text{HC}}^{\text{SCP}} \quad h_{\text{Sys}}^{\text{SCP}} \quad h_{\text{Cart}}^{\text{SCP}}}{h_{\text{HC, Sys, Cart}}^{\text{SCP}}} \quad \max(h_{\text{HC}}^{\text{SCP}}, h_{\text{Sys}}^{\text{SCP}}, h_{\text{Cart}}^{\text{SCP}})$$

Tasks (1667)	805	852	965	1006	1032
---------------------	-----	-----	-----	------	-------------

$$\max(\langle h_1 \rangle, \langle h'_1, h'_2 \rangle) \leq \max(\langle h_1, h'_1, h'_2 \rangle, \langle h'_1, h'_2, h_1 \rangle)$$

Drawbacks of Diversification

Diversification algorithm

- sample 1000 states
- start with empty set of orders
- until time limit is reached:
 - generate a random order
 - if it improves upon current set of orders, keep it
 - otherwise, discard it

Drawbacks of Diversification

Diversification algorithm

- sample 1000 states
 - start with empty set of orders
 - until time limit is reached:
 - generate a random order
 - if it improves upon current set of orders, keep it
 - otherwise, discard it
-
- considers only **random** orders
→ too many orders to stumble over good ones

Drawbacks of Hill Climbing

Hill climbing search

- sample 1000 states
- start with random order
- until no better successor for samples found:
 - swap positions of two heuristics
 - move to first improving successor

Drawbacks of Hill Climbing

Hill climbing search

- sample 1000 states
 - start with random order
 - until no better successor for samples found:
 - swap positions of two heuristics
 - move to first improving successor
-
- tries to find order for **set of states**
 - there may not be a single order for multiple states
 - starts with **random** order
 - decent initial solution important for local optimization

Improvements

- optimize for **single** state
- start with **greedy** order
- diversify **optimized** greedy orders

Greedy Order

Objectives:

- increase heuristic value quickly
- preserve costs for subsequent heuristics

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Value-per-cost ratio

$$\text{ratio}(h, s) = \frac{h(s)}{\text{saturated costs for } h}$$

Greedy Order

Objectives:

- increase heuristic value quickly
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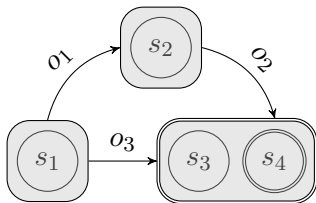
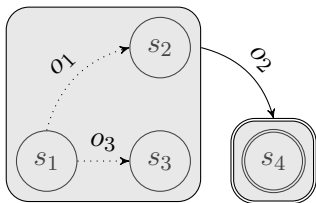
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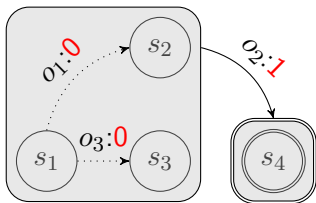
Greedy algorithm

- start with empty order
- until all heuristics are ordered
 - append heuristic h with highest value-per-cost ratio
 - subtract saturated costs for h from overall costs

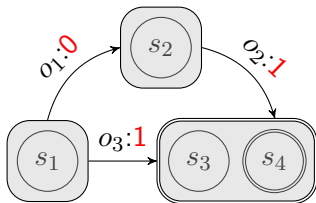
Greedy Order: Example



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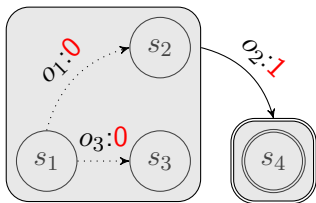


$h_1(s_1) = 1$
saturated costs: **1**
ratio: 1

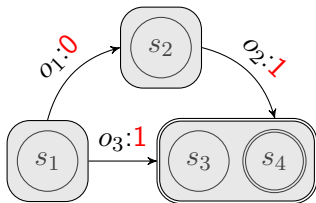


$h_2(s_1) = 1$
saturated costs: **2**
ratio = 0.5

Greedy Order: Example



$h_1(s_1) = 1$
saturated costs: **1**
ratio: 1



$h_2(s_1) = 1$
saturated costs: **2**
ratio = 0.5

$\rightarrow \langle h_1, h_2 \rangle$

Single Order for Initial State

	random	greedy	random-opt	greedy-opt
random	–	105	0	20
greedy	989	–	268	0
random-opt	1086	463	–	111
greedy-opt	1108	548	402	–

New diversification algorithm

Until time limit is reached:

- sample state s
- find **greedy** order for s
- **optimize** order with hill climbing
- keep order if **diverse**

Diverse Orders

New diversification algorithm

Until time limit is reached:

- sample state s
- find **greedy** order for s
- **optimize** order with hill climbing
- keep order if **diverse**

	max	random	random-opt	greedy-opt
Tasks <small>(1667)</small>	1032	1006	1009	1048

Comparison to State of the Art

Using h^2 mutexes:

- SymBA_2^* outperforms $h_{\text{greedy-opt}}^{\text{SCP}}$ in 11 domains
- $h_{\text{greedy-opt}}^{\text{SCP}}$ outperforms SymBA_2^* in 22 domains
- SymBA_2^* solves 1008 tasks
- $h_{\text{greedy-opt}}^{\text{SCP}}$ solves 1084 tasks

Related ICAPS Talk and Poster

Seipp, Keller & Helmert (ICAPS 2017)

A Comparison of Cost Partitioning Algorithms for Optimal Classical Planning

- theoretical and empirical comparison of cost partitioning algorithms
- saturated cost partitioning usually method of choice for hill climbing PDBs, systematic PDBs, Cartesian abstractions and landmark heuristics

Presented on Wednesday in the first ICAPS session “Optimal Planning”

Conclusion

- new greedy algorithm for finding orders
- pair with optimization and diversification
- combine heterogeneous abstraction heuristics